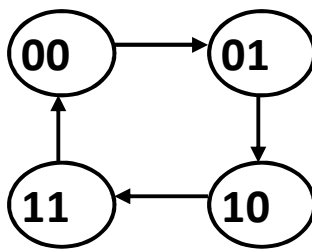


Sequential Circuit (Counter Design)

Synchronous (Parallel) Counters

- **Synchronous (parallel) counters:** the flip-flops are clocked at the same time by a common clock pulse.
- We can design these counters using the sequential logic design process (covered in Lecture #12).
- Example: 2-bit synchronous binary counter (using T flip-flops, or JK flip-flops with identical J,K inputs).



Present state		Next state		Flip-flop inputs	
A_1	A_0	A_1^+	A_0^+	TA_1	TA_0
0	0	0	1	0	1
0	1	1	0	1	1
1	0	1	1	0	1
1	1	0	0	1	1

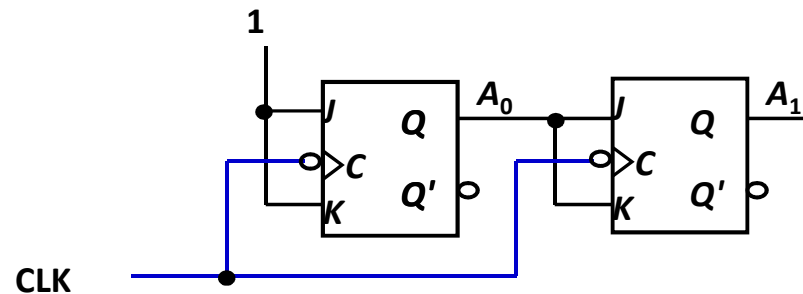
Synchronous (Parallel) Counters

- Example: 2-bit synchronous binary counter (using T flip-flops, or JK flip-flops with identical J,K inputs).

Present state		Next state		Flip-flop inputs	
A_1	A_0	A_1^+	A_0^+	TA_1	TA_0
0	0	0	1	0	1
0	1	1	0	1	1
1	0	1	1	0	1
1	1	0	0	1	1

$$TA_1 = A_0$$

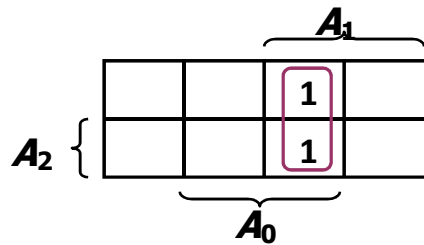
$$TA_0 = 1$$



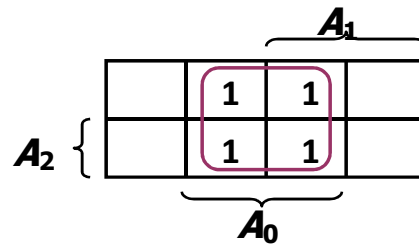
Synchronous (Parallel) Counters

- Example: 3-bit synchronous binary counter (using T flip-flops, or JK flip-flops with identical J, K inputs).

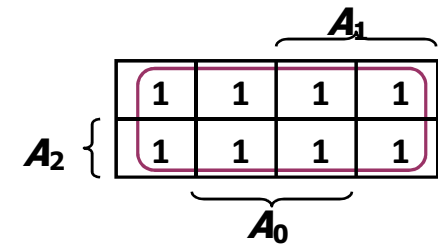
Present state			Next state			Flip-flop inputs		
A_2	A_1	A_0	A_2^+	A_1^+	A_0^+	TA_2	TA_1	TA_0
0	0	0	0	0	1	0	0	1
0	0	1	0	1	0	0	1	1
0	1	0	0	1	1	0	0	1
0	1	1	1	0	0	1	1	1
1	0	0	1	0	1	0	0	1
1	0	1	1	1	0	0	1	1
1	1	0	1	1	1	0	0	1
1	1	1	0	0	0	1	1	1



$$TA_2 = A_1 \cdot A_0$$



$$TA_1 = A_0$$



$$TA_0 = 1$$

Synchronous (Parallel) Counters

- Note that in a binary counter, the n^{th} bit (shown underlined) is always complemented whenever

$$\underline{0}11\dots11 \rightarrow \underline{1}00\dots00$$

$$\text{or } \underline{1}11\dots11 \rightarrow \underline{0}00\dots00$$

- Hence, X_n is complemented whenever $X_{n-1}X_{n-2} \dots X_1X_0 = 11\dots11$.
- As a result, if T flip-flops are used, then

$$TX_n = X_{n-1} \cdot X_{n-2} \cdot \dots \cdot X_1 \cdot X_0$$

Synchronous (Parallel) Counters

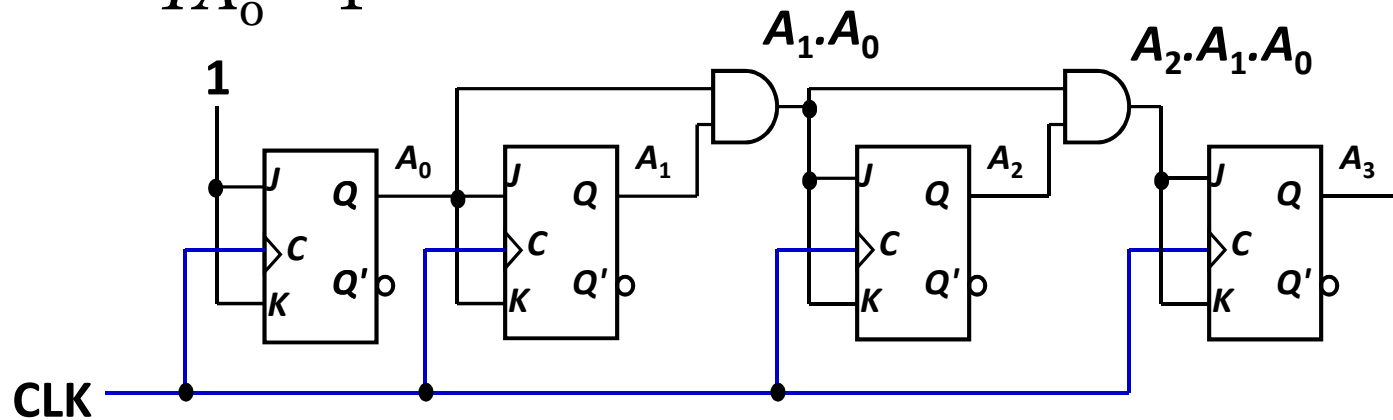
- Example: 4-bit synchronous binary counter.

$$TA_3 = A_2 \cdot A_1 \cdot A_0$$

$$TA_2 = A_1 \cdot A_0$$

$$TA_1 = A_0$$

$$TA_0 = 1$$



Synchronous (Parallel) Counters

- Example: Synchronous decade/BCD counter.

Clock pulse	Q_3	Q_2	Q_1	Q_0
Initially	0	0	0	0
1	0	0	0	1
2	0	0	1	0
3	0	0	1	1
4	0	1	0	0
5	0	1	0	1
6	0	1	1	0
7	0	1	1	1
8	1	0	0	0
9	1	0	0	1
10 (recycle)	0	0	0	0

$$T_0 = 1$$

$$T_1 = Q_3' \cdot Q_0$$

$$T_2 = Q_1 \cdot Q_0$$

$$T_3 = Q_2 \cdot Q_1 \cdot Q_0 + Q_3 \cdot Q_0$$

Synchronous (Parallel) Counters

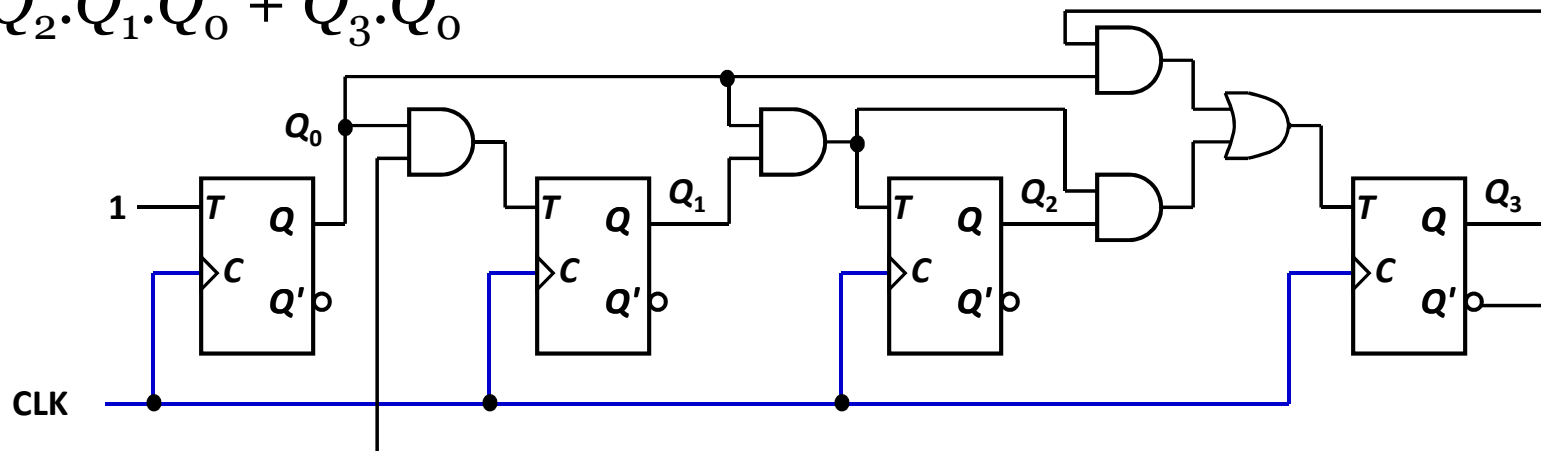
- Example: Synchronous decade/BCD counter (cont'd).

$$T_0 = 1$$

$$T_1 = Q_3' \cdot Q_0$$

$$T_2 = Q_1 \cdot Q_0$$

$$T_3 = Q_2 \cdot Q_1 \cdot Q_0 + Q_3 \cdot Q_0$$

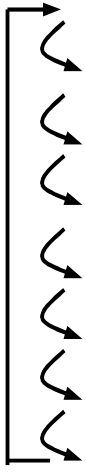



Up/Down Synchronous Counters

- **Up/down synchronous counter:** a *bidirectional* counter that is capable of counting either up or down.
- An input (control) line Up/\overline{Down} (or simply Up) specifies the direction of counting.
 - ❖ $Up/\overline{Down} = 1 \rightarrow$ Count upward
 - ❖ $Up/\overline{Down} = 0 \rightarrow$ Count downward

Up/Down Synchronous Counters

- Example: A 3-bit up/down synchronous binary counter.

Clockpulse	Up	Q_2	Q_1	Q_0	Down
0		0	0	0	
1		0	0	1	
2		0	1	0	
3		0	1	1	
4		1	0	0	
5		1	0	1	
6		1	1	0	
7		1	1	1	

Up counter

$$TQ_0 = 1$$

$$TQ_1 = Q_0$$

$$TQ_2 = Q_0 \cdot Q_1$$

Down

counter

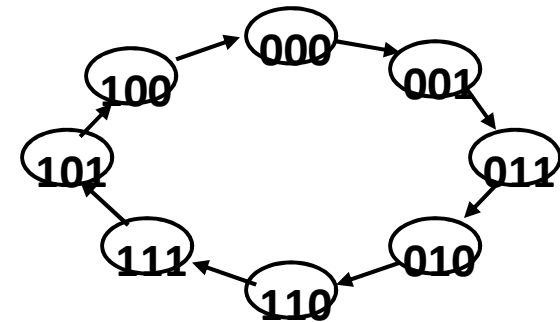
$$TQ_0 = 1$$

$$TQ_1 = Q_0'$$

$$TQ_2 = Q_0' \cdot Q_1'$$

Designing Synchronous Counters

- Covered in Lecture #12.
- Example: A 3-bit Gray code counter (using JK flip-flops).



Present state			Next state			Flip-flop inputs					
Q_2	Q_1	Q_0	Q_2^+	Q_1^+	Q_0^+	JQ_2	KQ_2	JQ_1	KQ_1	JQ_0	KQ_0
0	0	0	0	0	1	0	X	0	X	1	X
0	0	1	0	1	1	0	X	1	X	X	0
0	1	0	1	1	0	1	X	X	0	0	X
0	1	1	0	1	0	0	X	X	0	X	1
1	0	0	0	0	0	X	1	0	X	0	X
1	0	1	1	0	0	X	0	0	X	X	1
1	1	0	1	1	1	X	0	X	0	1	X
1	1	1	1	0	1	X	0	X	1	X	0

Cont...

- 3-bit Gray code counter: flip-flop inputs.

	Q_1	Q_0		
Q_2	00	01	11	10
0				1
1	X	X	X	X

$JQ_2 = Q_1 \cdot Q_0'$

	Q_1	Q_0		
Q_2	00	01	11	10
0	X	X	X	X
1	1			

$KQ_2 = Q_1' \cdot Q_0'$

	Q_1	Q_0		
Q_2	00	01	11	10
0		1	X	X
1			X	X

$JQ_1 = Q_2' \cdot Q_0$

	Q_1	Q_0		
Q_2	00	01	11	10
0	X	X		
1	X	X	1	

$KQ_1 = Q_2 \cdot Q_0$

	Q_1	Q_0		
Q_2	00	01	11	10
0	1	X	X	
1		X	X	1

$JQ_0 = Q_2 \cdot Q_1 + Q_2' \cdot Q_1'$
 $= (Q_2 \oplus Q_1)'$

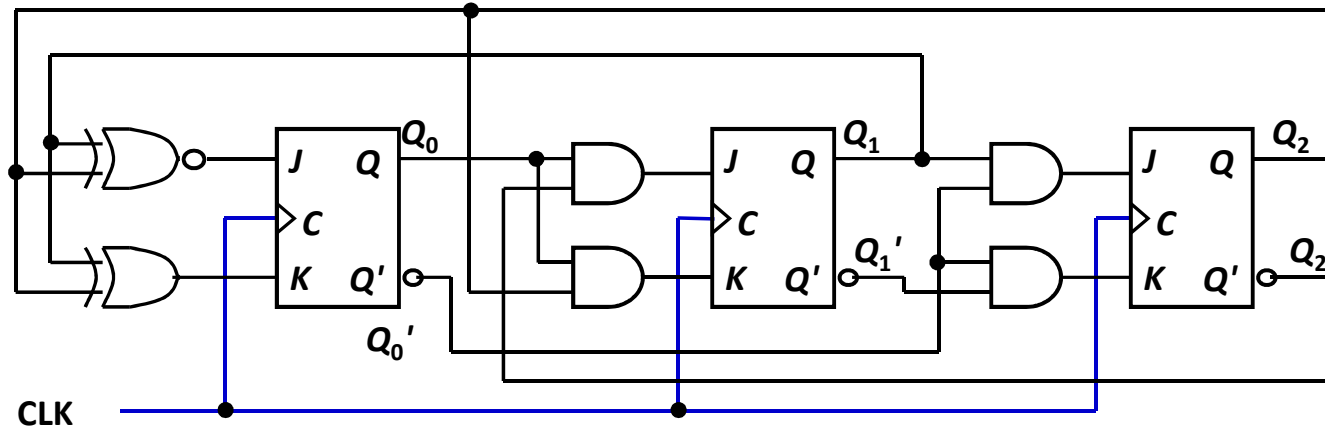
	Q_1	Q_0		
Q_2	00	01	11	10
0	X		1	X
1	X	1		X

$KQ_0 = Q_2 \cdot Q_1' + Q_2' \cdot Q_1$
 $= Q_2 \oplus Q_1$

Cont...

- 3-bit Gray code counter: logic diagram.

$$JQ_2 = Q_1 \cdot Q_0' \quad JQ_1 = Q_2' \cdot Q_0 \quad JQ_0 = (Q_2 \oplus Q_1)'$$
$$KQ_2 = Q_1' \cdot Q_0' \quad KQ_1 = Q_2 \cdot Q_0 \quad KQ_0 = Q_2 \oplus Q_1$$

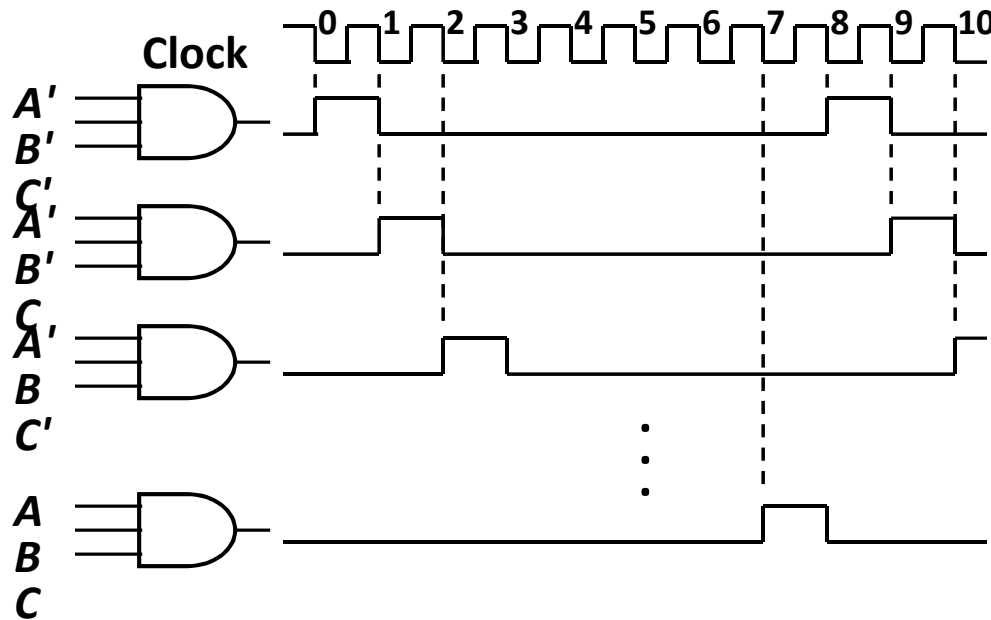


Decoding A Counter

- Decoding a counter involves determining which state in the sequence the counter is in.
- Differentiate between *active-HIGH* and *active-LOW* decoding.
- Active-HIGH decoding: output HIGH if the counter is in the state concerned.
- Active-LOW decoding: output LOW if the counter is in the state concerned.

Cont...

- Example: MOD-8 ripple counter (active-HIGH decoding).



**HIGH only on count
of $ABC = 000$**

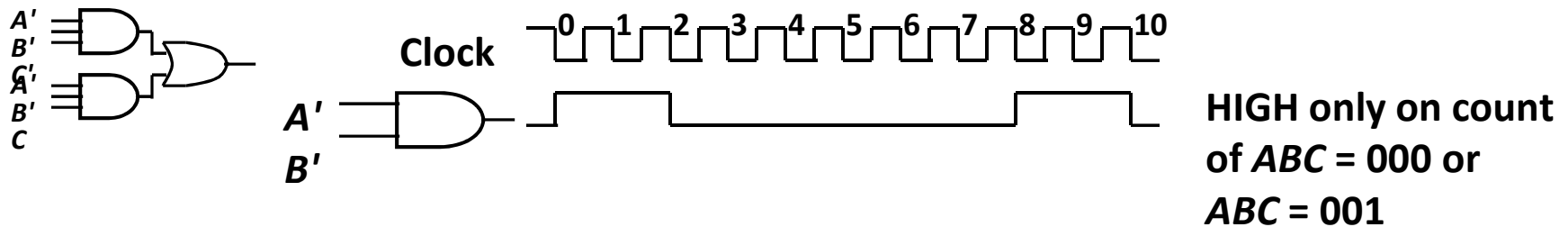
**HIGH only on
count of $ABC = 001$**

**HIGH only on
count of $ABC = 010$**

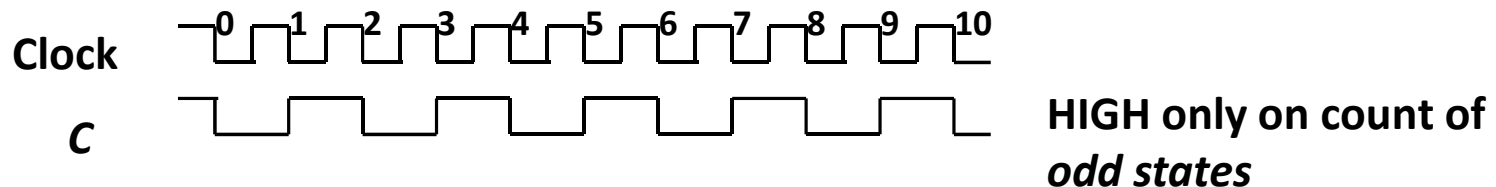
**HIGH only on
count of $ABC = 111$**

Cont...

- Example: To detect that a MOD-8 counter is in state 0 (000) or state 1 (001).



- Example: To detect that a MOD-8 counter is in the odd states (states 1, 3, 5 or 7), simply use C .

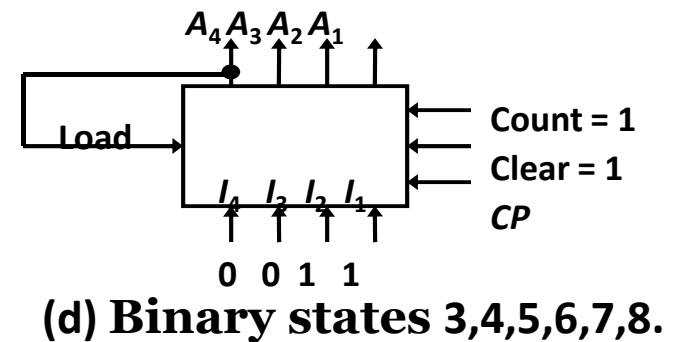
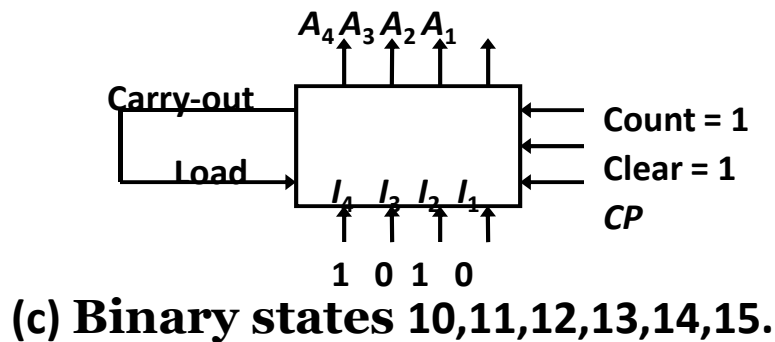
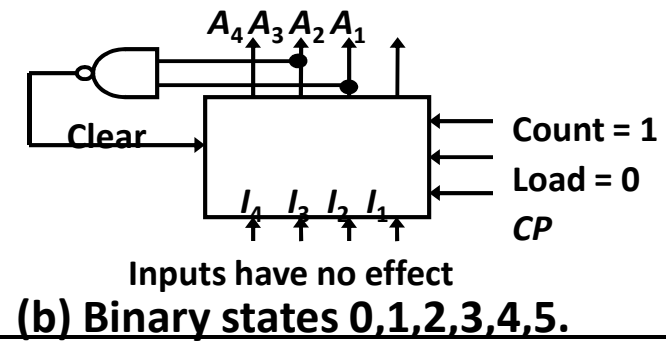
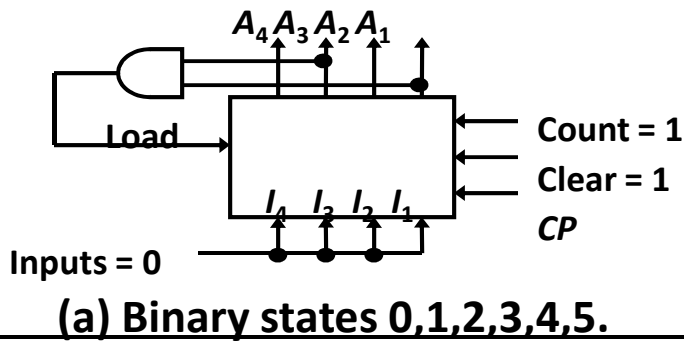


Counters with Parallel Load

- Counters could be augmented with parallel load capability for the following purposes:
 - ❖ To start at a different state
 - ❖ To count a different sequence
 - ❖ As more sophisticated register with increment/decrement functionality.

Counters with Parallel Load

- Different ways of getting a MOD-6 counter:



Counters with Parallel Load

- 4-bit counter with parallel load.

Clear	CP	Load	Count	Function
0	X	X	X	Clear to 0
1	X	0	0	No change
1	↑	1	X	Load inputs
1	↑	0	1	Next state

